CS 109
Spring 2008
Theory of Computation: Advanced

Homework 8
Due Wed May 7, 5:00pm

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General Instructions: Same as in Homework 1.

Honor Principle: Same as in Homework 1.

16. Prove that $\mathsf{NP} \subseteq \mathsf{BPP}$ implies $\mathsf{NP} = \mathsf{RP}$.

    Hint: Once you "solve" one NP-complete problem, you can solve them all! [2 points]

17. Let X and Y be finite sets and let Y^X denote the set of all functions from X to Y. We will think of these functions as "hash" functions.* A family $\mathcal{H} \subseteq Y^X$ is said to be 2-universal if the following property holds, with $h \in \mathcal{H}$ picked uniformly at random:

$$\forall\, x, x' \in X \;\forall\, y, y' \in Y \left(x \neq x' \;\Rightarrow\; \Pr_h[h(x) = y \wedge h(x') = y'] = \frac{1}{|Y|^2} \right) .$$

    Consider the sets $X = \{0,1\}^n$ and $Y = \{0,1\}^k$, with $k \leq n$. Treat the elements of X and Y as column vectors with 0/1 entries. For a matrix $A \in \{0,1\}^{k \times n}$ and vector $b \in \{0,1\}^k$, define the function $h_{A,b} : X \to Y$ as follows: $h_{A,b}(x) = Ax + b$, where all additions and multiplications are performed mod 2.

    Prove that $\{h_{A,b} : A \in \{0,1\}^{k \times n}, b \in \{0,1\}^k\}$ is a 2-universal family of hash functions. [2 points]

*The term "hash function" has no formal meaning; instead, one should speak of a "family of hash functions" or a "hash family" as we do here.