## Let your Mach-O fly

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**Black Hat Briefings** 

#### Who am I?

- Student at Politecnico di Milano.
- Security Consultant at Secure Network srl.
- Reverse Engineer at Zynamics GmbH.



#### Goal of the talk

In-memory execution of arbitrary binaries on a Mac OS X machine.

#### Talk outline

- Mach-O file structure
- XNU binary execution
- Attack technique
- Defeat ASLR on libraries to enhance the attack

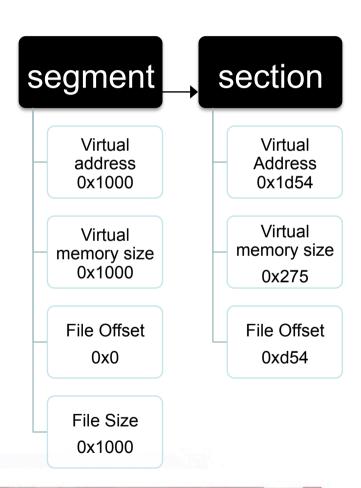
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#### Mach-O file

- Header structure: information on the target architecture and options to interpret the file.
- Load commands: symbol table location, registers state.
- Segments: define region of the virtual memory, contain sections with code or data.

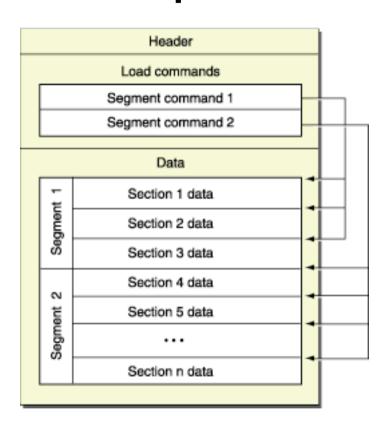
#### Segment and Sections



## Important segments

- PAGEZERO, if a piece of code accesses
   NULL it lands here. no protection flags.
- TEXT, holds code and read-only data. RX protection.
- DATA, holds data. RW protection.
- LINKEDIT, holds information for the dynamic linker including symbol and string tables. RW protection.

# Mach-O representation



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## Binary execution

- Conducted by the kernel and the dynamic linker.
- The kernel, when finishes his part, jumps to the dynamic linker entry point.
- The dynamic linker is not randomized.

## **Execution steps**

#### Kernel

- Maps the dynamic linker in the process address space.
- Parses the header structure and loads all segments.
- Creates a new stack.

#### **Dynamic linker**

- Retrieves base address of the binary.
- Resolves symbols.
- Resolves library dependencies.
- Jumps to the binary entry point.

#### Stack

- Mach-O file base address.
- Command line arguments.
- Environment variables.
- Execution path.
- All padded.

## Stack representation

Mach-o
Address
Argc
Argv[]

0
Envp[]

0
exec\_path ptr

0
exec\_path
Argv[] strings
Envp[] strings

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## Proposed attack

- Userland-exec attack.
- Encapsulate a shellcode, aka autoloader, and a crafted stack in the injected binary.
- Execute the auto-loader in the address space of the attacked process.

#### WWW

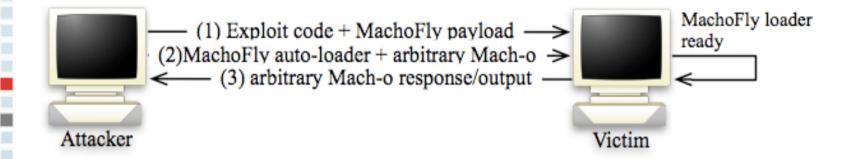
- Who: an attacker with a remote code execution in his pocket.
- Where: the attack is two-staged. First run a shellcode to receive the binary, then run the auto-loader contained in the binary.
  - Why: later in this talk.

#### What kind of binaries?

Any Mach-O file, from Is to Safari

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## A nice picture



## Infected binary

- We need to find a place to store the auto-loader and the crafted stack.
- PAGEZERO infection technique.
  - Cavity infector technique.

#### PAGEZERO INFECTION

- Change \_\_PAGEZERO protection flags with a custom value.
- Store the crafted stack and the autoloader code at the end of the binary.
- Point \_\_PAGEZERO to the crafted stack.
- Overwrite the first bytes of the file with the auto-loader address.

## Binary layout

**MODIFIED HEADER** INFECTED \_\_PAGEZERO load commands and segments sections and binary data SHELLCODE **CRAFTED STACK** 

#### Auto-loader

- Impersonates the kernel.
- Un-maps the old binary.
- Maps the new one.

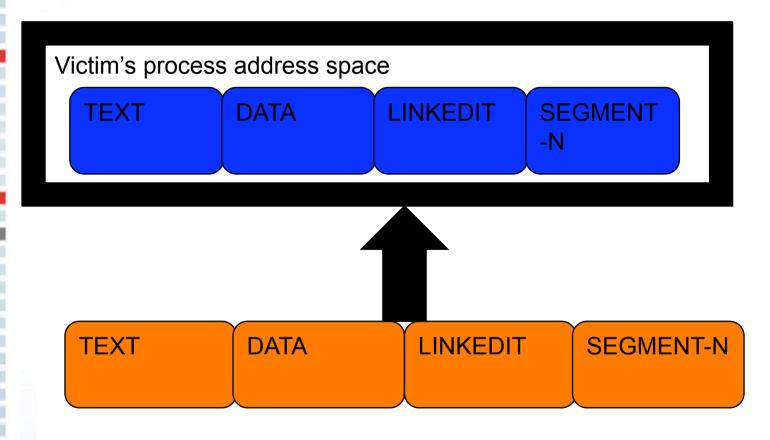
## Auto-loader description

- Parses the binary.
- Reads the virtual addresses of the injected binary segments.
- Unloads the attacked binary segments pointed by the virtual addresses.
- Loads the injected binary segments.

## Auto-loader description(2)

- Maps the crafted stack referenced by PAGEZERO.
- Cleans registers.
- Cleans some libSystem variables.
- Jumps to dynamic linker entry point.

## We do like pictures, don't we?



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## libSystem variables

- \_malloc\_def\_zone\_state
- \_NXArgv\_pointer
- \_malloc\_num\_zones
- \_\_keymgr\_global

# Why are those variables important?

- They are used in the initialization of malloc.
- Two of them are used for command line arguments parsing.
- Not cleaning them will result in a crash.

#### Hunts the variables

- Mac OS X Leopard has ASLR for libraries.
- Those variables are not exported.
- Cannot use dlopen()/dlsym() combo.

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#### **Defeat ASLR**

- Retrieve libSystem in-memory base address.
- Read symbols from the libSystem binary.
  - Adjust symbols to the new address.

## How ASLR works in Leopard

- Only libraries are randomized.
- The randomization is performed whenever the system or the libraries are updated.
- Library segments addresses are saved in dyld\_shared\_cache\_arch.map.

#### Retrieve libSystem address

- Parse
   dyld\_shared\_cache\_
   i386.map and
   search for libSystem
   entry.
- Adopt functions exported by the dynamic linker and perform the whole task in-memory.

#### **Dyld functions**

- \_dyld\_image\_count() used to retrieve the number of linked libraries of a process.
- \_dyld\_get\_image\_header() used to retrieve the base address of each library.
- \_dyld\_get\_image\_name() used to retrieve the name of a given library.

#### Find 'em

- Parse dyld load commands.
- Retrieve \_\_LINKEDIT address.
- Iterate dyld symbol table and search for the functions name in \_\_LINKEDIT.

## Back to libSystem

- Non-exported symbols are taken out from the symbol table when loaded.
- Open libSystem binary, find the variables in the symbol table.
- Adjust variables to the base address of the in-memory \_\_\_DATA segment.

## Put pieces together

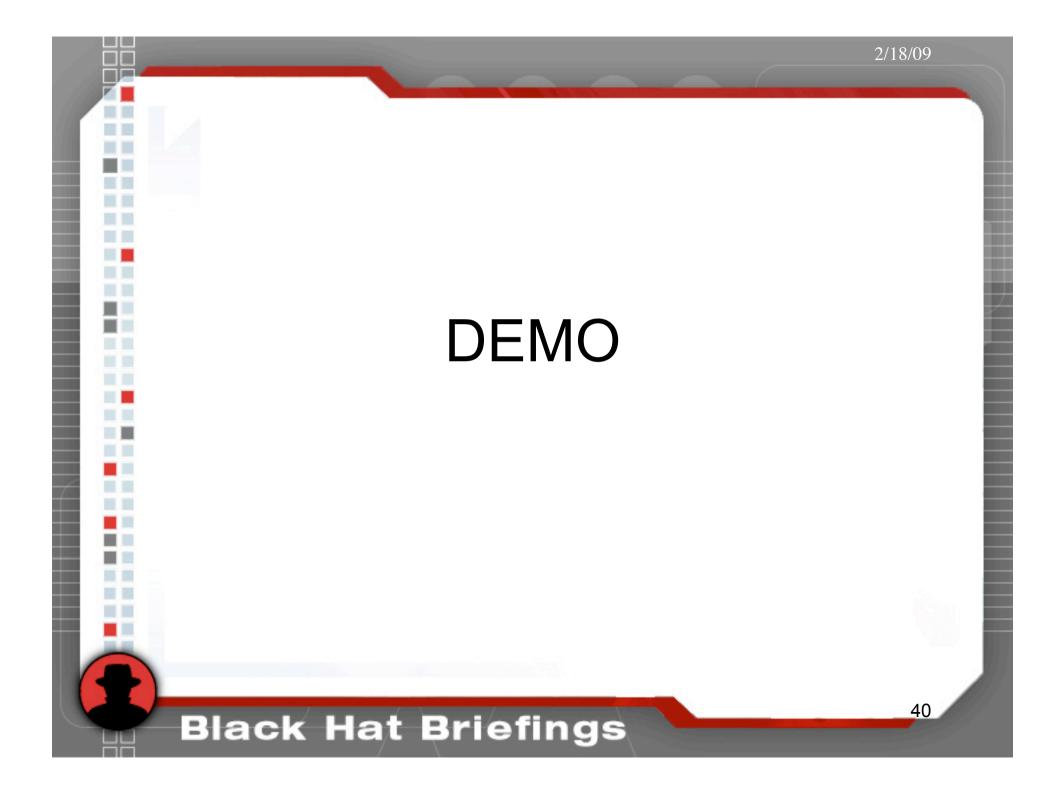
- Iterate the header structure of libSystem in-memory and find the \_\_\_DATA base address.
  - DATA base address 0x2000
  - Symbol at 0x2054
  - In-memory \_\_\_DATA base address 0x4000
  - Symbol in-memory at 0x4054

#### Results

- Run a binary into an arbitrary machine.
- No traces on the hard-disk.
- No execve(), the kernel doesn't know about us.
- It works with every binary.
- It is possible to write payloads in a high level language.

## Demo description

- Run a simple piece of code which acts like a shellcode and retrieve the binary.
- Execute the attack with nmap and Safari.
- Show network dump.
- Show memory layout before and after the attack.



## Future developments

- Employ encryption to avoid NIDS detection.
- Using cavity infector technique.
- Port the code to iPhone to evade code signing protection (Catch you at BH Europe).

## Thanks, questions?