Homework 1
Due Wed Apr 9, 5:00pm

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General Instructions: Please write concisely, but rigorously. Non-rigorous solutions won't be graded. For each problem, only "nearly flawless" solutions earn 2 points. Solutions that contain the key insights but are flawed in execution earn only 1 point. The purpose of this strict grading scheme is to dissuade you from writing up half-baked ideas in the hope of getting "some" credit.

Honor Principle: You are allowed to discuss the problems and exchange solution ideas with your classmates. But when you write up any solutions for submission, you must work alone. You may refer to any textbook you like, including online ones. However, you may not refer to published or online solutions to the specific problems on the homework, if you intend to turn it in for credit. *If in doubt, ask the professor for clarification!*

1. For a complexity class $\mathcal{C}$, define two new complexity classes "$\exists\mathcal{C}$" and "$\forall\mathcal{C}$" as follows.

$$\begin{aligned}\exists\mathcal{C} &= \{\{x \in \Sigma^* : \exists y \in \Sigma^* \text{ with } |y| = \text{poly}(|x|) \text{ such that } \langle x, y\rangle \in L_0\} : L_0 \in \mathcal{C}\} \\ \forall\mathcal{C} &= \{\{x \in \Sigma^* : \forall y \in \Sigma^* \text{ with } |y| = \text{poly}(|x|) \text{ we have } \langle x, y\rangle \in L_0\} : L_0 \in \mathcal{C}\}\end{aligned}$$

The notation $|y| = \text{poly}(|x|)$ means that there is some fixed polynomial p such that $|y| = O(p(|x|))$.

Prove, rigorously, that $\exists\mathsf{P} = \mathsf{NP}$ and $\forall\mathsf{P} = \mathsf{coNP}$. Use only the basic definitions from class, where NP was defined using NDTMs and coNP was defined as $\{L \subseteq \Sigma^* : \overline{L} \in \mathsf{NP}\}$. (This problem is simply asking you to write out a rigorous version of ideas we have already discussed in class.) [2 points]

2. Define the following two complexity classes

$$\begin{aligned}\mathsf{EXPTIME} &= \bigcup_{i=1}^{\infty} \mathsf{DTIME}(2^{n^i}) \\ \mathsf{NEXPTIME} &= \bigcup_{i=1}^{\infty} \mathsf{NTIME}(2^{n^i}) .\end{aligned}$$

Prove that $\mathsf{P} = \mathsf{NP}$ implies $\mathsf{EXPTIME} = \mathsf{NEXPTIME}$. The key trick is to "pad" an input with a lot of extra symbols. [2 points]

3. For this problem, assume that Boolean formulas are encoded as strings over the alphabet $\{0, 1, \vee, \wedge, \neg, (,)\}$, fully parenthesized to resolve ambiguities. Note that the negation operator $(\neg)$ has higher priority than the other two. The variables in a formula ϕ are represented as binary substrings of ϕ with no leading zeros. For instance, the formula

$$((x_1 \wedge \neg x_2) \vee \neg x_1 \vee (x_3 \wedge \neg x_1)) \wedge \neg x_4 \wedge (\neg x_3 \vee x_5)$$

is represented as the string

$$((1 \wedge \neg 10) \vee \neg 1 \vee (11 \wedge \neg 1)) \wedge \neg 100 \wedge (\neg 11 \vee 101) .$$

Define $\textsc{satisfies} = \{\langle \phi, \alpha\rangle : \phi \text{ is a Boolean formula and the assignment } \alpha \text{ satisfies } \phi\}$. Our proof that $\textsc{sat} \in \mathsf{NP}$ boiled down to showing that $\textsc{satisfies} \in \mathsf{P}$. Prove the stronger result that $\textsc{satisfies} \in \mathsf{L}$ (i.e., LOGSPACE). Recall, from class, that $\mathsf{L} \subseteq \mathsf{P}$.

This problem is all about careful implementation, so take care to specify exactly how you use the work tape of your TM. Some naïve implementations end up requiring $\Omega(\log^2 n)$ space. [2 points]