

**General Instructions.** Same as for Homework 1.

**Honor Principle.** Same as for Homework 1.

---

4. The *extreme points problem* asks whether the convex hull of  $n$  given points in the plane has  $n$  vertices (i.e., whether all of the  $n$  points are “extreme”); note that this is potentially an easier problem than actually computing the convex hull.

Model this problem as a set recognition problem, i.e., that of recognizing whether or not an input vector  $\mathbf{x}$  belongs to  $W$ , for an appropriate set  $W \subseteq \mathbb{R}^{2n}$ . Prove that the number of connected components  $\#W \geq (n-1)!$  and conclude that the algebraic computation tree complexity of the problem is  $\Omega(n \log n)$ .

5. Let  $\mathbf{a}_1, \dots, \mathbf{a}_k$  and  $\mathbf{b}$  be fixed nonzero vectors in  $\mathbb{R}^n$  such that the system of inequalities

$$\langle \mathbf{a}_1, \mathbf{x} \rangle \geq 0, \langle \mathbf{a}_2, \mathbf{x} \rangle \geq 0, \dots, \langle \mathbf{a}_k, \mathbf{x} \rangle \geq 0,$$

in the unknown  $\mathbf{x} \in \mathbb{R}^n$ , is feasible and implies the inequality  $\langle \mathbf{b}, \mathbf{x} \rangle \geq 0$ . Here  $\langle \cdot, \cdot \rangle$  denotes the standard inner product in  $\mathbb{R}^n$ . Then it can be shown that  $\mathbf{b}$  is a non-negative linear combination of the  $\mathbf{a}_i$ 's, i.e.,  $\mathbf{b} = \sum_{i=1}^k \lambda_i \mathbf{a}_i$  for some non-negative reals  $\{\lambda_i\}$ . This fact is sometimes known as Farkas's Lemma.

Using Farkas's Lemma, prove the following two lower bounds in the *linear* decision tree model (i.e., on input  $\mathbf{x}$ , each internal node gets to ask a question “ $\sum_{i=1}^n c_i x_i \geq 0$ ?” where the  $c_i$ 's are constants).

- 5.1. The complexity of finding the largest of  $n$  given reals is  $n-1$ .  
5.2. The complexity of finding the second largest is at least  $n-2 + \log n$ .

Hint: Once you have solved #5.1, use what you learnt along with a leaf counting argument to solve #5.2.