

# CS 30: Discrete Math in CS (Winter 2020): Lecture 5

Date: 13th January, 2020 (Monday)

Topic: Proofs via Contradiction

*Disclaimer: These notes have not gone through scrutiny and in all probability contain errors.*

*Please discuss in Piazza/email errors to [deeparnab@dartmouth.edu](mailto:deeparnab@dartmouth.edu)*

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This is one of the most commonly used styles of proof. When faced with a proposition  $p$  (either in propositional logic, or predicate logic – often the latter) which we wish to prove true, we *suppose for the sake of contradiction* that  $p$  were false. Then we logically deduce something *absurd* (like  $0 = 1$  or  $3$  is even), that is, something which we know to be false. This implies that our supposition (which is,  $p$  is false) must be wrong. Therefore, the proposition  $p$  must be true. This method of proving is also called *reductio ad absurdum* — reduction to absurdity.

Formally, in the jargon of logic, what the above argument captures is the fact that the following formula

$$(\neg p \Rightarrow \text{false}) \Rightarrow p$$

is a *tautology*. Can you deduce this from the equivalences?

A final word before we move on to concrete examples. Many times the false is obtained by showing that some other proposition  $q$  holds as well as its negation. That is, we end up showing  $(\neg p \Rightarrow (q \wedge \neg q))$ . Interestingly, sometimes this proposition is  $p$  itself.

Just this lecture, we write down the steps in a list so as to make sure all ideas are clear.

## 1. A Simple Warm-up.

**Lemma 1.** For all numbers  $n$ , if  $n^2$  is even, then  $n$  is even.

*Proof.*

- (a) Suppose, for the sake of contradiction, the proposition is *not true*.
- (b) That is, there exists a number  $n$  such that  $n^2$  is even but  $n$  is not even. That is,  $n$  is odd. (We negated the predicate logic statement).
- (c) Since  $n$  is odd,  $n = 2k + 1$  for some integer  $k$ .
- (d) This implies  $n^2 = (2k + 1)^2 = 4k^2 + 4k + 1 = 2(2k^2 + 2k) + 1$ .
- (e) That is,  $n^2$  is odd. This is a contradiction to  $P(n^2)$ , that is,  $n^2$  is even.
- (f) Therefore, our supposition must be wrong, that is, the proposition is true.

□



**Exercise:** Mimic the above proof to prove: For any number  $n$ , if  $n^2$  is divisible by 3, then  $n$  is divisible by 3.



**Exercise:** Prove by contradiction: the product of a non-zero rational number and an irrational number is irrational.

## 2. A Pythagorean<sup>1</sup> Theorem.

**Theorem 1.**  $\sqrt{2}$  is irrational.

*Proof.*

- (a) Suppose, for the sake of contradiction, that  $\sqrt{2}$  is indeed rational.
- (b) Since  $\sqrt{2}$  is rational, there exists two integers  $a, b$  such that  $\sqrt{2} = a/b$ .
- (c) By dividing out common factors, we may assume  $\gcd(a, b) = 1$ .
- (d) Since  $a/b = \sqrt{2}$ , we get  $a = \sqrt{2} \cdot b$ . Squaring both sides, we get  $a^2 = 2b^2$ .
- (e) Therefore  $a^2$  is even.
- (f) Lemma 1 implies that  $a$  is even. And therefore  $a = 2\ell$  for some  $\ell$ .
- (g) Therefore,  $a^2 = 4\ell$ .
- (h) Since  $a^2 = 2b^2$ , we get  $4\ell = 2b^2$ , which in turn implies  $b^2 = 2k$ . That is,  $b^2$  is even.
- (i) Lemma 1 implies that  $b$  is even.
- (j) Thus, we have deduced both  $a$  and  $b$  are even. This **contradicts**  $\gcd(a, b) = 1$ .
- (k) Therefore, our supposition that  $\sqrt{2}$  is rational must be wrong. That is,  $\sqrt{2}$  is irrational.

□



**Exercise:** Mimic the above proof to prove that  $\sqrt{3}$  is irrational. How far can you generalize? Can you prove that  $\sqrt{n}$  is irrational if  $n$  is not a perfect square, that is,  $n$  is not  $a^2$  for some integer  $a$ ?

## 3. A Euclidean Theorem. Here is another classic example of Proof by Contradiction.

**Theorem 2.** There are infinitely many primes.

*Proof.*

- (a) Suppose, for the sake of contradiction, there were only finitely many primes.
- (b) Let  $q$  be the largest of these primes.
- (c) Therefore, for any number  $n > q$ ,  $n$  is *not* a prime.
- (d) Consider the number  $n = q! + 1$ . Recall,  $q! = 1 \times 2 \times \dots \times q$ .

<sup>1</sup>This is of course not the famous Pythagorean theorem on right angled triangles, but nonetheless a Pythagorean may be the first to have proved it. See [https://en.wikipedia.org/wiki/Irrational\\_number](https://en.wikipedia.org/wiki/Irrational_number), for instance.

- (e) Since  $n > q$ , this  $n$  is not a prime.
- (f) Therefore, there exists some prime  $p$  such that  $p \mid n$ .
- (g) Since  $q$  is the largest prime,  $p \leq q$ .
- (h) But this means  $p \mid q!$ , which means  $p \nmid q! + 1$ . That is,  $p \nmid n$ .
- (i) We have deduced both  $p \mid n$  and  $p \nmid n$ . Contradiction. Thus our supposition is wrong. There are infinitely many primes.

□

#### 4. The AM-GM inequality

**Theorem 3.** If  $a$  and  $b$  are two positive real numbers, then  $a + b \geq 2\sqrt{ab}$ .

*Proof.*

- (a) Suppose for the sake of contradiction, there existed positive reals  $a, b$  with  $a + b < 2\sqrt{ab}$ .
- (b) Since both sides of the above inequality are positive, we can square both sides. That is,  $(a + b)^2 < (2\sqrt{ab})^2$ .  
*Please note how crucial the fact that both sides were positive is. Otherwise, we cannot square and maintain the inequality. And indeed, the theorem is incorrect for negative numbers. Consider  $a = -1$  and  $b = -1$ . The RHS is 2 but the LHS is  $-2$ .*
- (c) That is,  $a^2 + 2ab + b^2 < 4ab$ .
- (d) That is,  $a^2 - 2ab + b^2 < 0$ .
- (e) That is,  $(a - b)^2 < 0$ .
- (f) But  $(a - b)^2 \geq 0$ , since it is a square. Thus, we have reached a contradiction.

□