CS 10: Problem solving via Object Oriented Programming Winter 2017

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Day 16 – Graph Traversals





2. Breadth first search

Graph traversals are useful to answer questions about vertex relationships

Some Graph traversals uses

- Uses are typically around *reachability*
- Computing *path* from vertex u to vertex v
- Given start vertex s of Graph G, compute a path with the minimum number of edges between s and all other vertices (or report no such path exists)
- Testing whether G is fully connected (e.g., all vertices reachable)
- Identifying *cycles* in G (or reporting no cycle exists)
- Today's examples have no cycles (next class will consider them)

Depth First Search (DFS) uses a stack to explore as if in a maze



DFS basic idea

- Keep going until you can't go any further, then back track
- Relies on a stack (implicit or explicit) to keep track of where you've been

Some of you did Depth First Search on Problem Set 1

RegionFinder

Loop over all the pixels If a pixel is unvisited and of the correct color Start a new region Keep track of pixels need to be visited, initially just one As long as there's some pixel that needs to be visited Get one to visit Add it to the region Mark it as visited Loop over all its neighbors If the neighbor is of the correct color Add it to the list of pixels to be visited If the region is big enough to be worth keeping, do so

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If you added to end of list...

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Loop over all the pixels If a pixel is unvisited and of the correct color Start a new region Keep track of pixels need to be visited, initially just one As long as there's some pixel that needs to be visited Get one to visit And if you get pixel from end of Add it to the region list, you implemented a stack Mark it as visited Loop over all its neighbors If the neighbor is of the correct color Add it to the list of pixels to be visited If the region is big enough to be worth keeping, do so

If you added to end of list...



```
stack.push(s) //start node
repeat until find goal vertex or
stack empty:
    u = stack.pop()
    if !u.visited
        u.visited = true
        (maybe do something while here)
        for v ∈ u.adjacent
            if !v.visited
            stack.push(v)
```



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            if !v.visited
            stack.push(v)
```



DFS algorithm

stack.push(s) //start node repeat until find goal vertex or stack empty: u = stack.pop()if !u.visited u.visited = true(maybe do something while here) for $v \in u.adjacent$ if !v.visited stack.push(v)

Push unvisited adjacent



```
stack.push(s) //start node
repeat until find goal vertex or
stack empty:
    u = stack.pop()
    if !u.visited
        u.visited = true
        (maybe do something while here)
        for v ∈ u.adjacent
            if !v.visited
            stack.push(v)
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DFS algorithm

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        (maybe do something while here)
        for v ∈ u.adjacent
            if !v.visited
            stack.push(v)
```

Push unvisited adjacent (F, but not A)

\sum



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stack.push(s) //start node
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DFS algorithm

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stack empty:
 u = stack.pop()
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 stack.push(v)

Push unvisited adjacent (H, but not B)

 \square



```
stack.push(s) //start node
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        for v ∈ u.adjacent
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DFS algorithm

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stack.push(s) //start node
repeat until find goal vertex or
stack empty:
    u = stack.pop()
    if !u.visited
        u.visited = true
        (maybe do something while here)
        for v ∈ u.adjacent
            if !v.visited
            stack.push(v)
```

Nothing to push, back up by popping C



```
stack.push(s) //start node
repeat until find goal vertex or
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        (maybe do something while here)
        for v ∈ u.adjacent
            if !v.visited
            stack.push(v)
```



DFS algorithm

stack.push(s) //start node
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stack empty:
 u = stack.pop()
 if !u.visited
 u.visited = true
 (maybe do something while here)
 for v ∈ u.adjacent
 if !v.visited
 stack.push(v)

Nothing to push, back up by popping D

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```
stack.push(s) //start node
repeat until find goal vertex or
stack empty:
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    if !u.visited
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        (maybe do something while here)
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DFS algorithm

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repeat until find goal vertex or
stack empty:
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 for v ∈ u.adjacent
 if !v.visited
 stack.push(v)

Push unvisited adjacent (G, but not A)



```
stack.push(s) //start node
repeat until find goal vertex or
stack empty:
    u = stack.pop()
    if !u.visited
        u.visited = true
        (maybe do something while here)
        for v ∈ u.adjacent
            if !v.visited
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   if !u.visited
       u.visited = true
        (maybe do something while here)
       for v ∈ u.adjacent
              if !v.visited
               stack.push(v)
```



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stack empty:
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        u.visited = true
        (maybe do something while here)
        for v ∈ u.adjacent
            if !v.visited
            stack.push(v)
```

Node discovery tells us something about the graph

Discovery edges

- Edges that lead to unvisited nodes
- Discovery edges form a tree on the graph
- Can traverse from start to goal on tree (if goal reachable)
- Can tell us which nodes are not reachable (not on path formed by discovery nodes)
- Not guaranteed to be shortest path!

Back edges

- Edges that lead to previously discovered nodes
- Lead to ancestor nodes in tree
- Indicate presence of a cycle in the graph

Run time is O(n+m)

Run time

- Assume graph with *n* nodes and *m* edges
- Visit each node at most one time (due to visited indicator)
- Visit each edge at most one time
- Run time is O(n+m)



- 1. Depth first search
- 2. Breadth first search

Breadth First Search (BFS) can find the shortest path between nodes



BFS basic idea

- Explore outward in "ripples"
- Look at all nodes 1 step away, then all nodes 2 steps away...
- Relies on a queue (implicit or explicit) implementation
- Path found from s to any other vertex is shortest

Some of you did Breadth First Search on Problem Set 1

RegionFinder

Loop over all the pixels If a pixel is unvisited and of the correct color Start a new region Keep track of pixels need to be visited, initially just one As long as there's some pixel that needs to be visited Get one to visit Add it to the region Mark it as visited Loop over all its neighbors If the neighbor is of the correct color Add it to the list of pixels to be visited If the region is big enough to be worth keeping, do so

Some of you did Breadth First Search on Problem Set 1

RegionFinder

Loop over all the pixels If a pixel is unvisited and of the correct color Start a new region Keep track of pixels need to be visited, initially just one As long as there's some pixel that needs to be visited Get one to visit Add it to the region Mark it as visited Loop over all its neighbors If the neighbor is of the correct color Add it to the list of pixels to be visited If the region is big enough to be worth keeping, do so

If you added to end of list...

Some of you did Breadth First Search on Problem Set 1

RegionFinder

Loop over all the pixels If a pixel is unvisited and of the correct color Start a new region Keep track of pixels need to be visited, initially just one As long as there's some pixel that needs to be visited Get one to visit And if you get pixel from front of Add it to the region list, you implemented a queue Mark it as visited Loop over all its neighbors If the neighbor is of the correct color Add it to the list of pixels to be visited If the region is big enough to be worth keeping, do so

If you added to end of list...

Breadth First Search (BFS) can find the shortest path between nodes



```
enqueue(s) //start node
s.visited = true
repeat until find goal vertex or
queue empty:
```

```
u = dequeque()
for v ∈ u.adjacent
    if !v.visited
        v.visited = true
        enqueue(v)
```

Breadth First Search (BFS) can find the shortest path between nodes



```
enqueue(s) //start node
s.visited = true
repeat until find goal vertex or
queue empty:
```

```
u = dequeque()
for v ∈ u.adjacent
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        v.visited = true
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enqueue(s) //start node
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repeat until find goal vertex or
queue empty:
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```
u = dequeque()
for v ∈ u.adjacent
    if !v.visited
        v.visited = true
        enqueue(v)
```



BFS algorithm

```
enqueue(s) //start node
s.visited = true
repeat until find goal vertex or
queue empty:
   u = dequeque()
```

```
for v ∈ u.adjacent
   if !v.visited
```

```
v.visited = true
enqueue (v)
```

dequeue(A), unvisited enqueue adjacent



BFS algorithm

```
enqueue(s) //start node
s.visited = true
repeat until find goal vertex or
queue empty:
   u = dequeque()
```

```
for v ∈ u.adjacent
   if !v.visited
```

```
v.visited = true
enqueue (v)
```

dequeue(B), enqueue unvisited adjacent F





BFS algorithm

```
enqueue(s) //start node
s.visited = true
repeat until find goal vertex or
queue empty:
    u = dequeque()
```

```
for v ∈ u.adjacent
if !v.visited
```

```
v.visited = true
enqueue(v)
```

dequeue(C), enqueue unvisited adjacent (none)





BFS algorithm

```
enqueue(s) //start node
s.visited = true
repeat until find goal vertex or
queue empty:
   u = dequeque()
```

enqueue (v)

v.visited = true

for v ∈ u.adjacent

if !v.visited

```
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```



```
enqueue(s) //start node
s.visited = true
repeat until find goal vertex or
queue empty:
```

```
u = dequeque()
for v ∈ u.adjacent
    if !v.visited
        v.visited = true
        enqueue(v)
```



```
enqueue(s) //start node
s.visited = true
repeat until find goal vertex or
queue empty:
    u = dequeque()
```

```
for v ∈ u.adjacent
    if !v.visited
        v.visited = true
```

```
enqueue (v)
```



v.visited = true

enqueue (v)









```
u = dequeque()
for v ∈ u.adjacent
    if !v.visited
        v.visited = true
        enqueue(v)
```

Node discovery tells us something about the graph

Discovery edges

- Lead to unvisited nodes
- Form a tree on the graph
- Can traverse from start to goal (or any node)
- Can tell us which nodes are not reachable (not on path formed by discovery nodes)
- Path guaranteed to have smallest number of edges

Can track how we got to node to find shortest path

- Build vertex tree
- Parent of each vertex is vertex that discovered it
- Parent is unique because we don't visit vertices twice

Run time is O(n+m)

Run time

- Assume graph with *n* nodes and *m* edges
- Visit each node at most one time (due to visited indicator
- Visit each edge at most one time
- Run time O(n+m)
- Useful for the Kevin Bacon game!